

Alternative Synchronization Strategies — Lock-Free Algorithms (1) —

lecture 19 (2021-05-24)

Master in Computer Science and Engineering

— Concurrency and Parallelism / 2020-21 —

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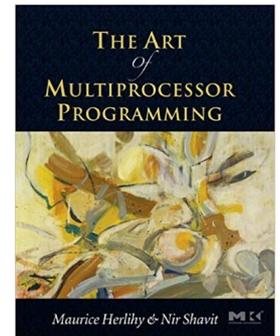
Alternative Synchronization Strategies

- Contents:

- Liveness: Types of Progress
- Coarse-Grained Synchronization
- Fine-Grained Synchronization
- Optimistic Synchronization
- Lazy Synchronization
- Lock-Free Synchronization

Past lectures

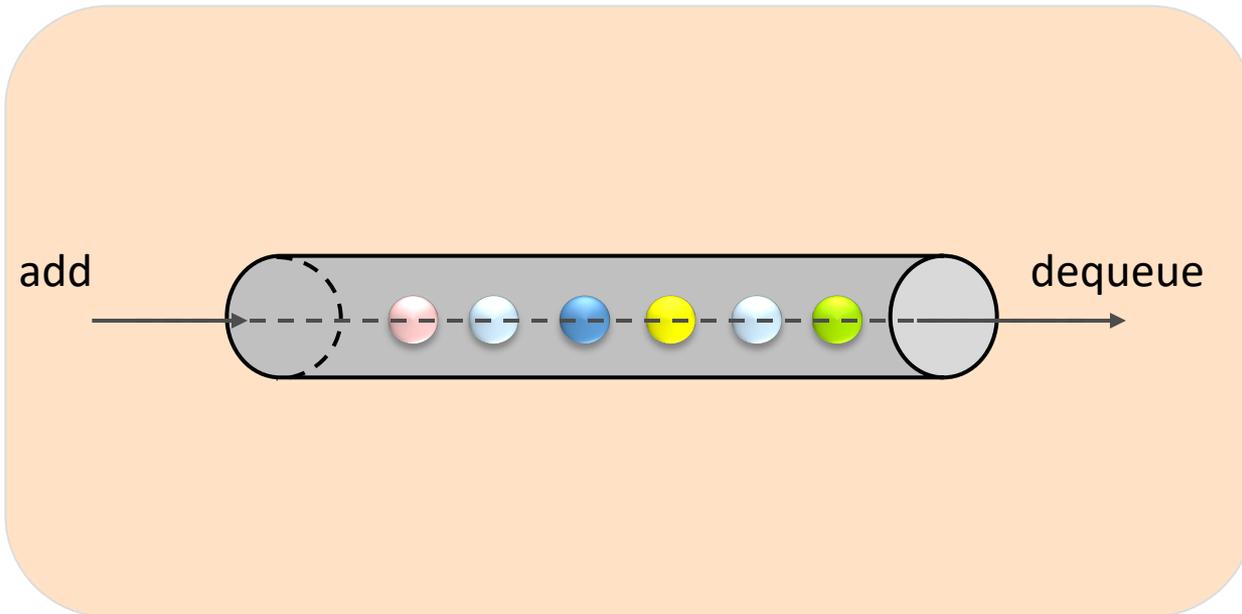
Today



- Reading list:

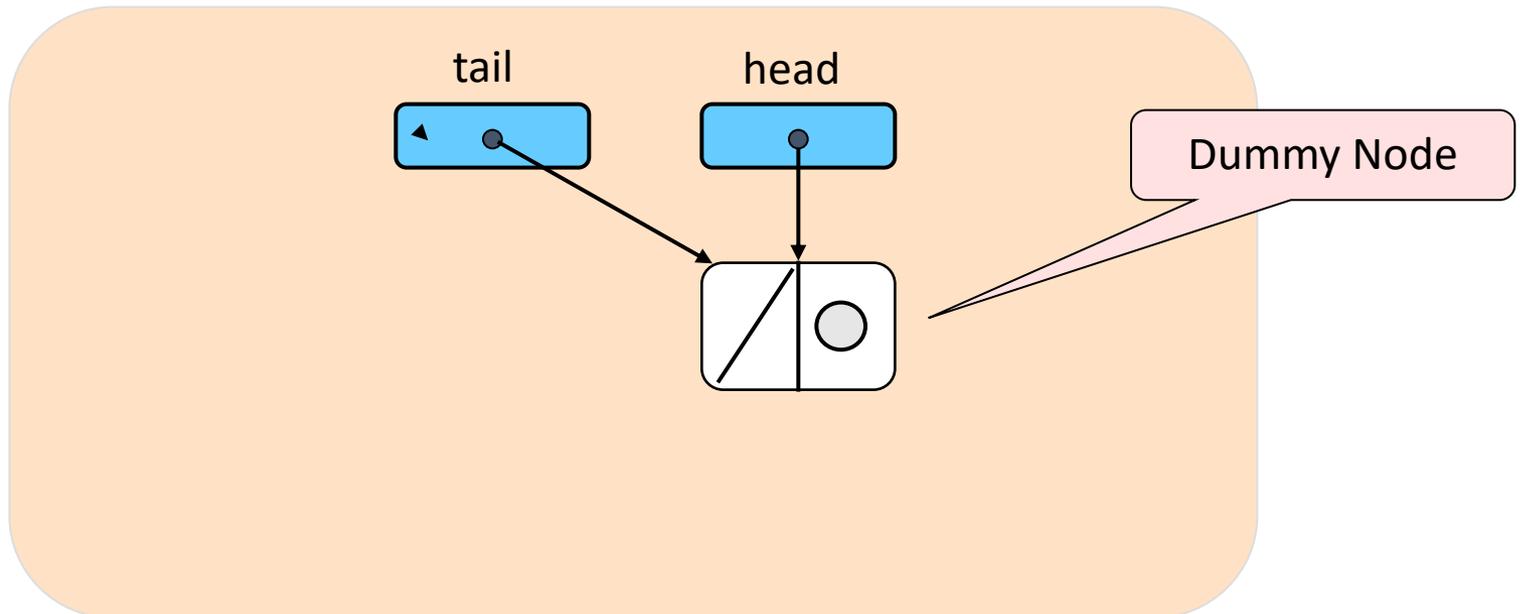
- chapter 5 of the Textbook
- Chapter 9 of “The Art of Multiprocessor Programming” by Maurice Herlihy & Nir Shavit (*available at clip*)

Basics for a lock-free Queue

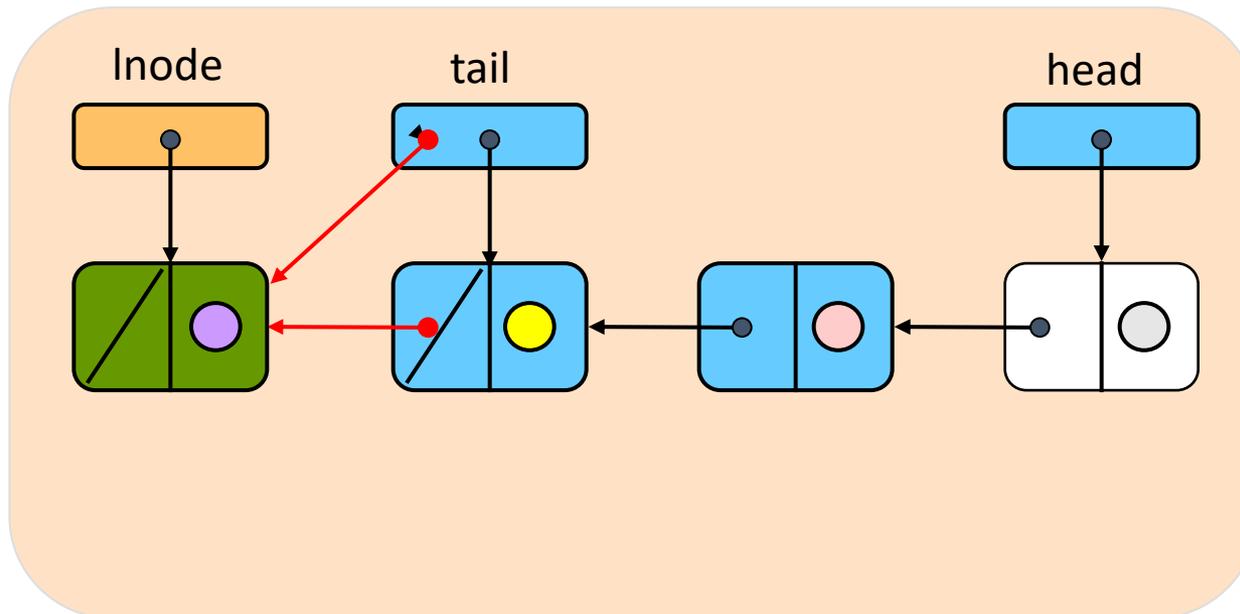


Basics for a lock-free Queue

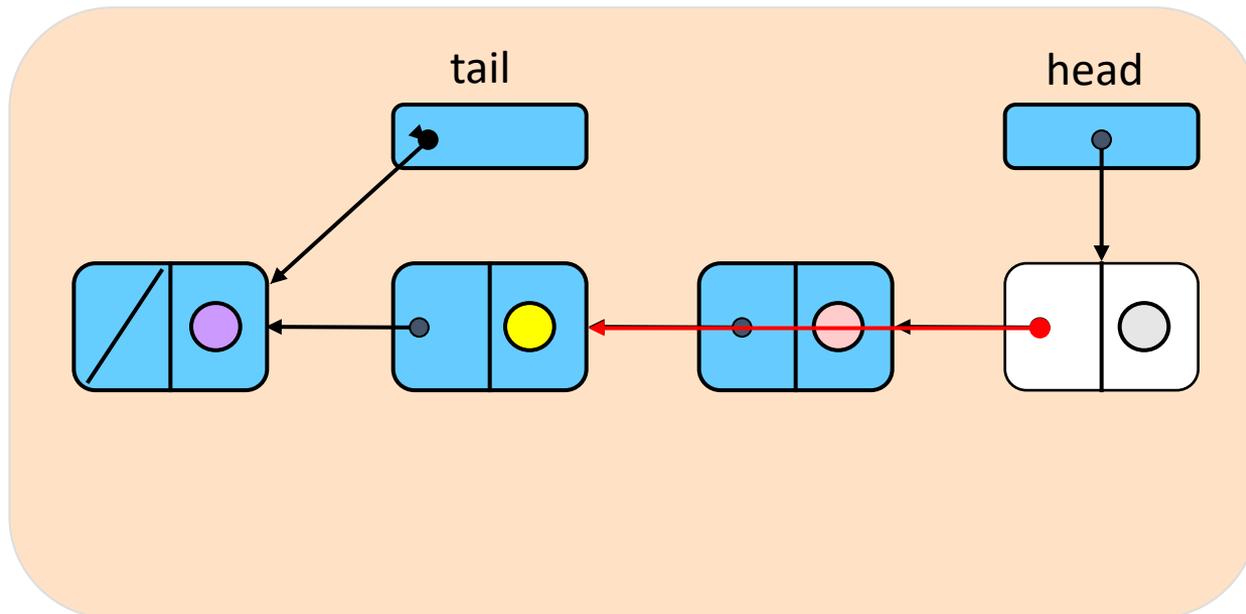
Empty queue



Enqueue



Deque



Compare & set (CAS)

shared
register old new
 ↓ ↓ ↓
CAS (A, B, C)

```
if A==B then A←C; return(true)
else return(false)
```

Supported by Intel, AMD, Arm, ...

Reminder: Lock-Free Data Structures

- No matter what ...



- Guarantees minimal progress in any execution
 - i.e., some thread will always complete a method call
- Even if others halt at malicious times
- Implies that implementation can't use locks

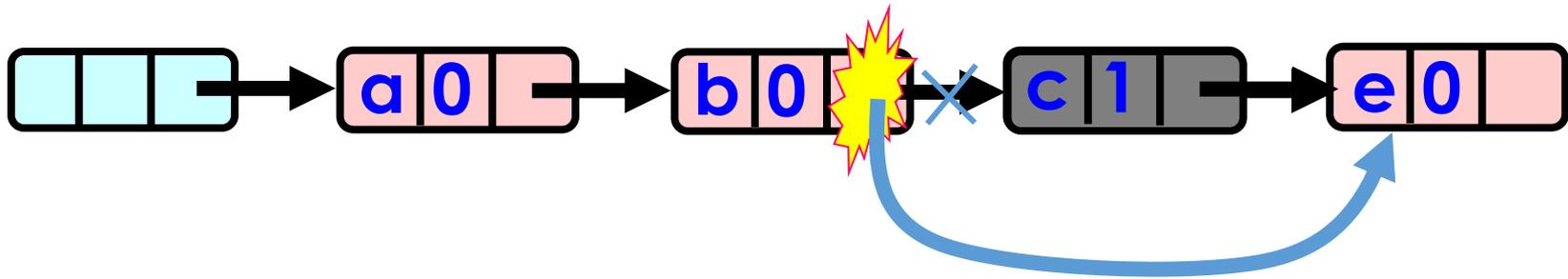
Lock-free Lists

- Next logical step (*after the lazy list*) is...
- Eliminate locking entirely
 - *contains()* wait-free
 - *add()* lock-free
 - *remove()* lock-free
- Use only *compareAndSet()*
- What could go wrong?

Remove Using CAS

- `remove(c)`

Logical Removal =
Set Mark Bit



Use CAS to verify
if pointer is correct

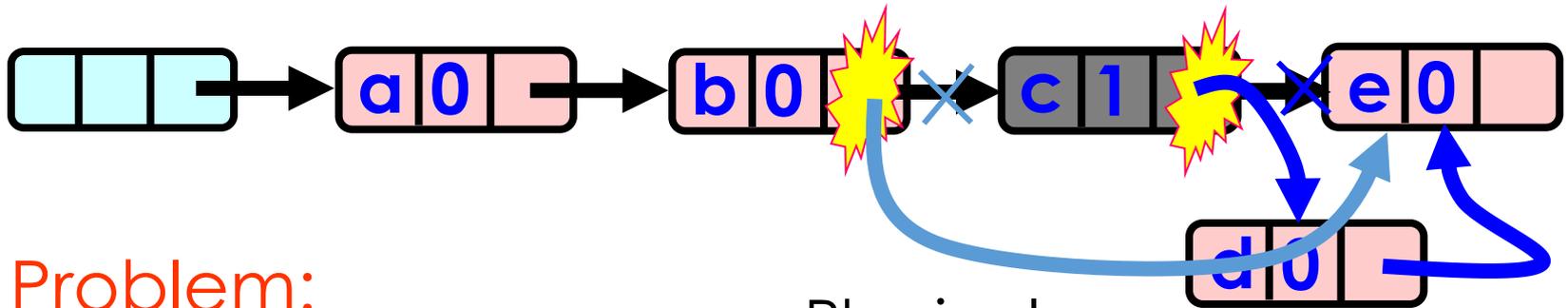
Physical
Removal
CAS pointer

Not enough! Why?

Problem...

- `remove(c) | add(d)`

Logical Removal =
Set Mark Bit



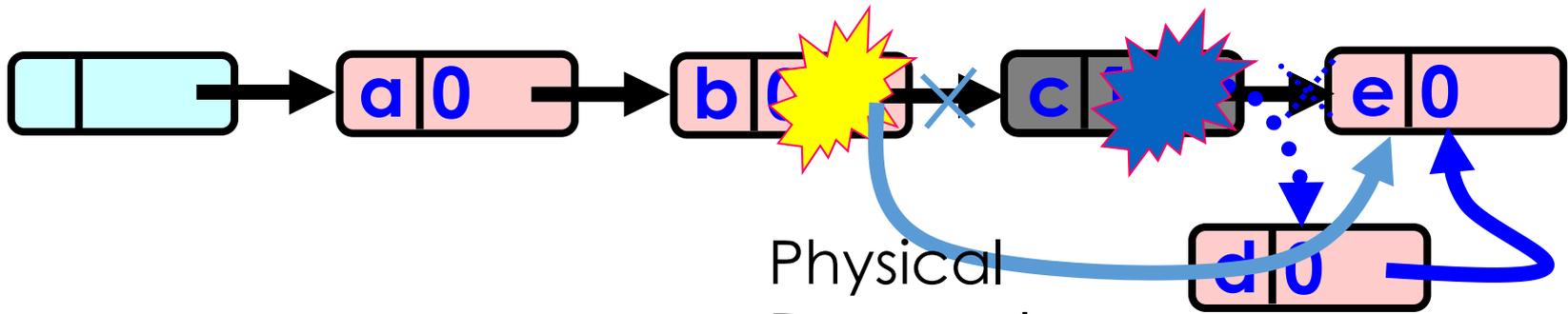
Problem:
'd' not added to list...
Must Prevent
manipulation of
removed node's pointer

Physical
Removal
CAS

Node added
Before
Physical
Removal CAS

The Solution: Combine Bit and Pointer

Logical Removal =
Set Mark Bit



Mark-Bit and Pointer
are CASed together
(AtomicMarkableReference)

Physical
Removal
CAS

Fail CAS: Node not
added after logical
Removal

Solution

- Use *AtomicMarkableReference*
- Atomically
 - Swing reference and
 - Update flag
- Remove in two steps
 - Set mark bit in next field
 - Redirect predecessor's pointer with a CAS

Marking a Node

- AtomicMarkableReference class
 - java.util.concurrent.atomic package



Extracting Reference & Mark

```
public Object get(boolean[] marked);
```

Extracting Reference & Mark

```
public object get(boolean[] marked);
```

**Returns
reference**

**Returns mark at
array index 0!**

Extracting Reference Only

```
public boolean isMarked();
```

Value of
mark

Changing State

```
public boolean compareAndSet(  
    Object expectedRef,  
    Object updateRef,  
    boolean expectedMark,  
    boolean updateMark);
```

Changing State

□

If this is the current
reference ...

```
public boolean compareAndSet(  
    Object expectedRef,  
    Object updateRef,  
    boolean expectedMark,  
    boolean updateMark);
```

And this is the
current mark ...

Changing State

□

...then change to this
new reference ...

```
public boolean compareAndSet(  
    Object expectedRef,  
    Object updateRef,  
    boolean expectedMark,  
    boolean updateMark);
```

... and this new
mark

Changing State

```
public boolean attemptMark(  
    Object expectedRef,  
    boolean updateMark);
```

Changing State

```
public boolean attemptMark(  
    Object expectedRef,  
    boolean updateMark);
```

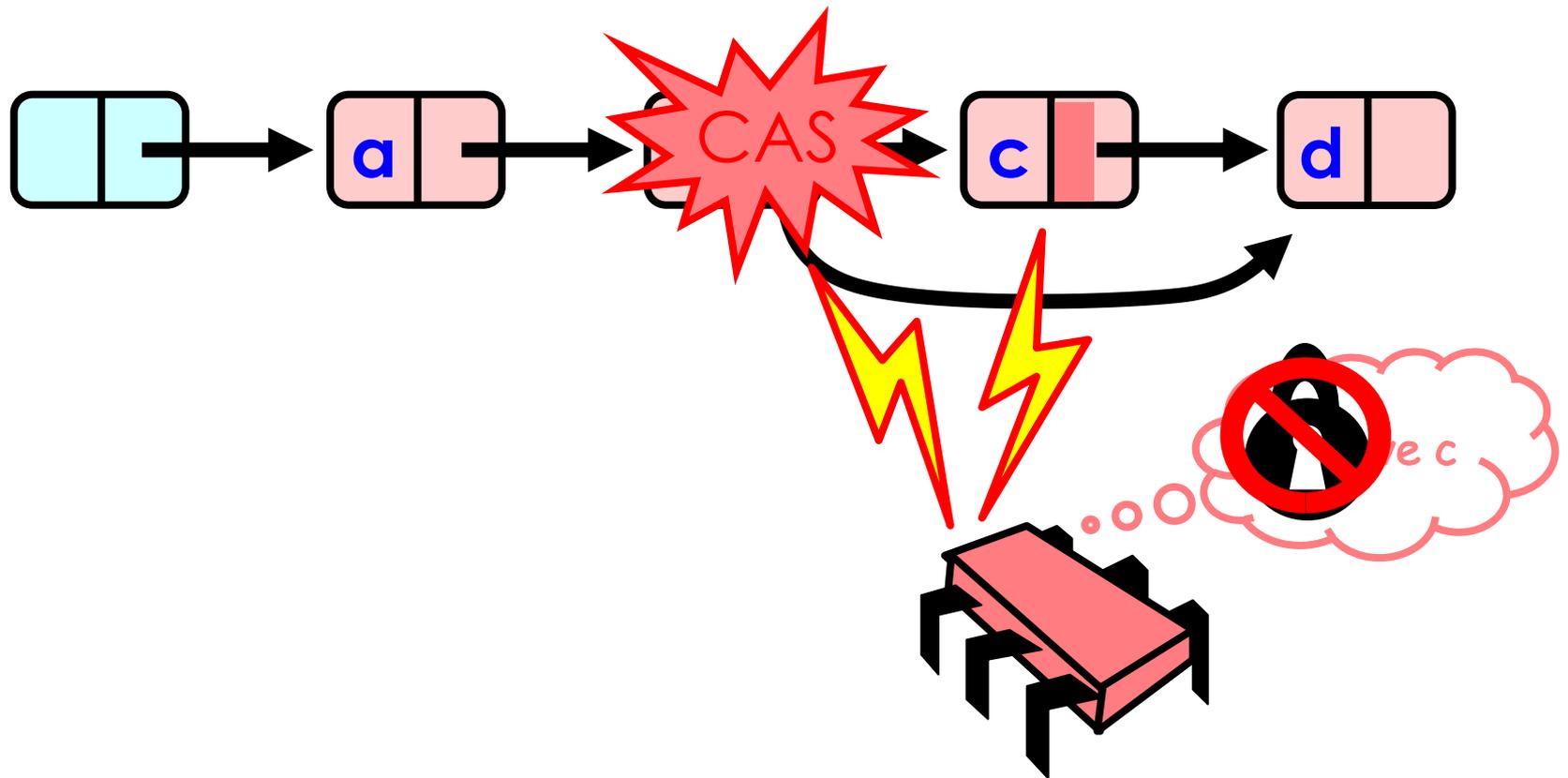
**If this is the current
reference ...**

Changing State

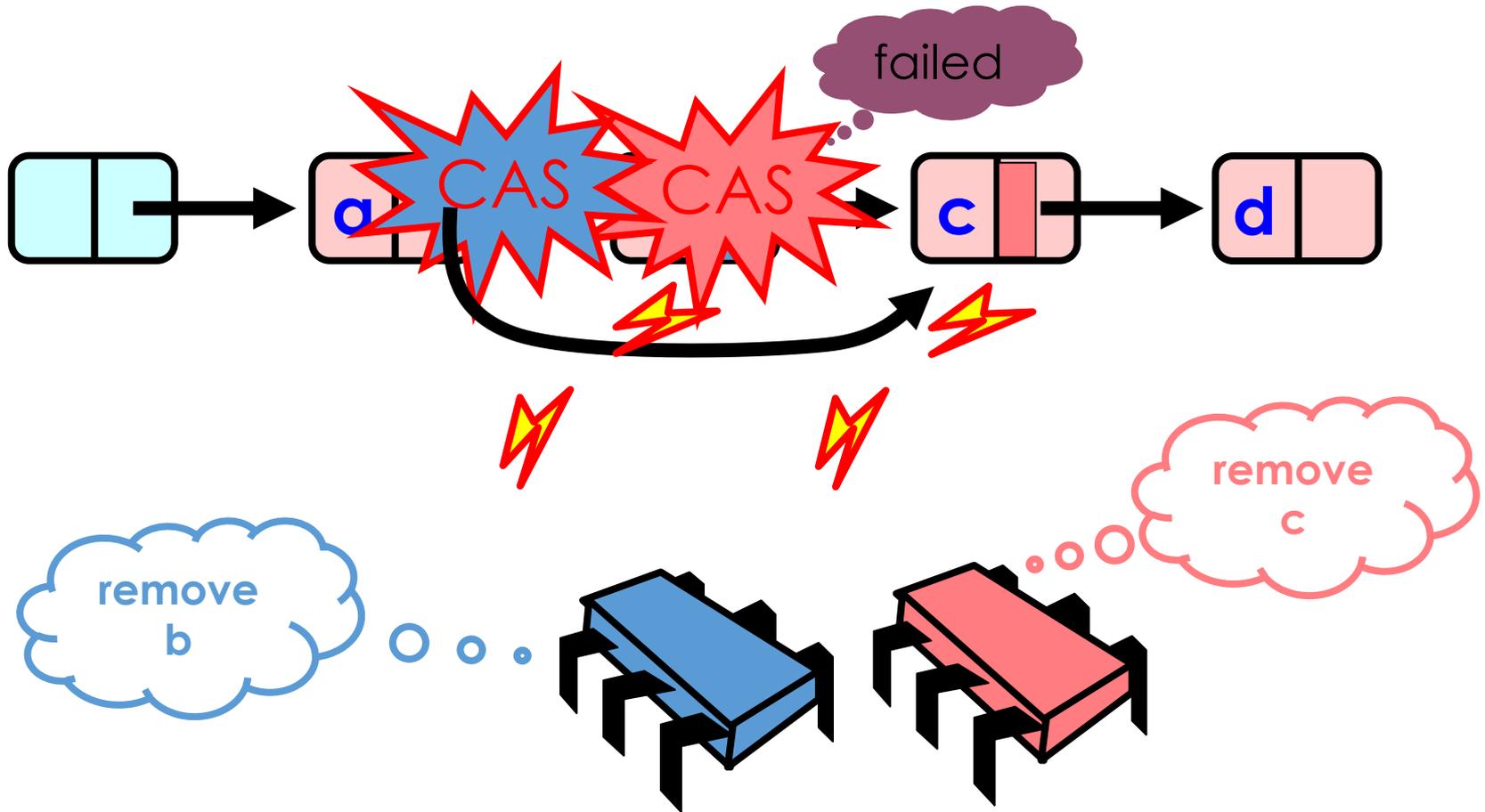
```
public boolean attemptMark(  
    Object expectedRef,  
    boolean updateMark);
```

**.. then change to
this new mark.**

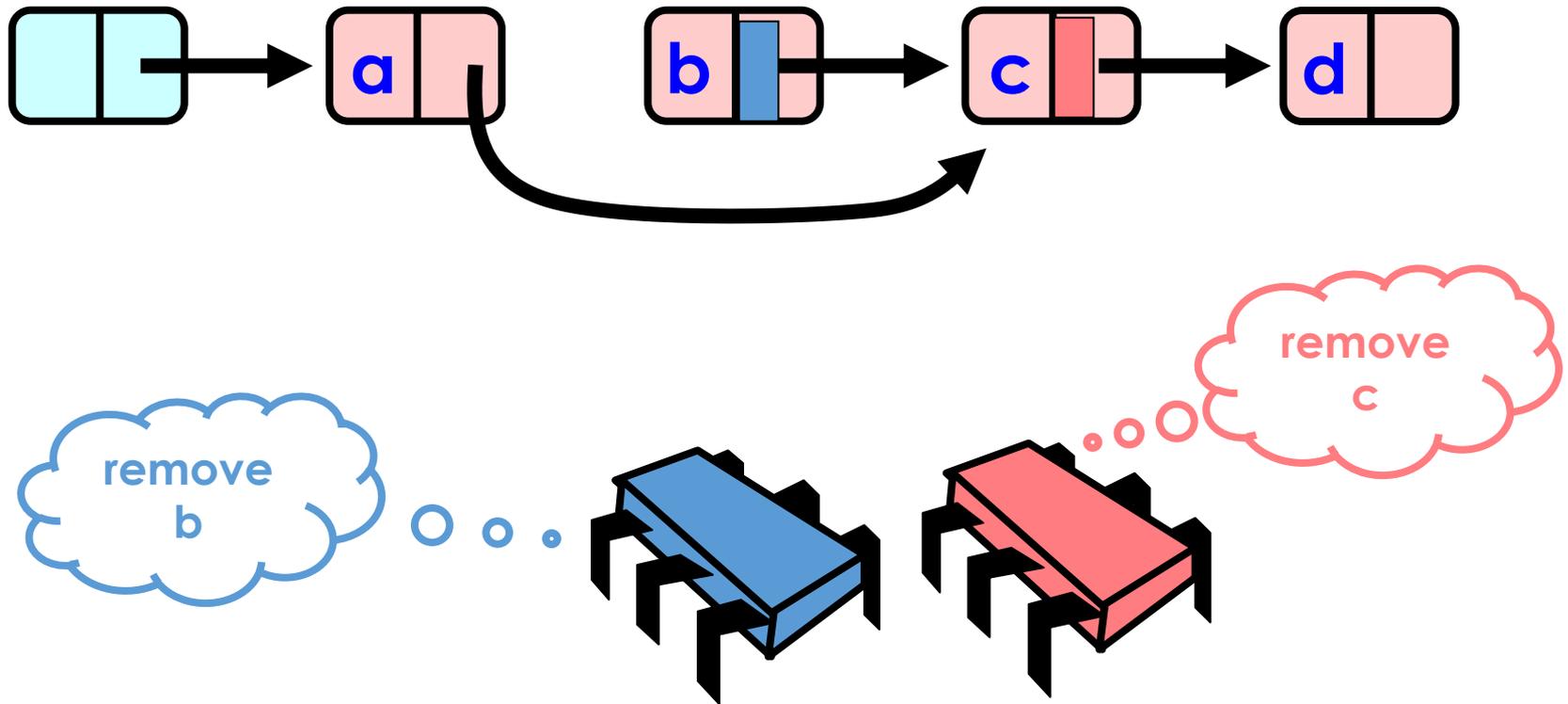
Removing a Node



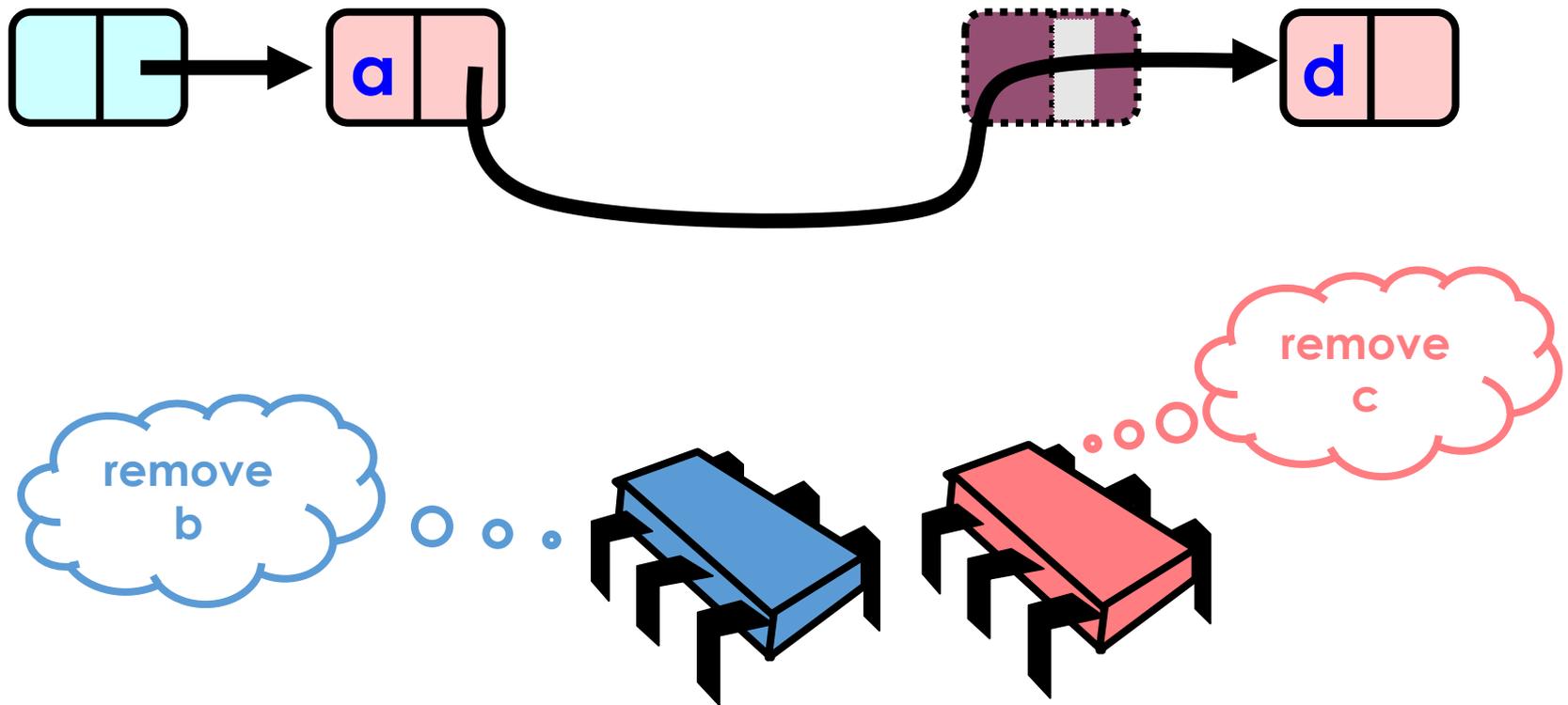
Removing a Node



Removing a Node



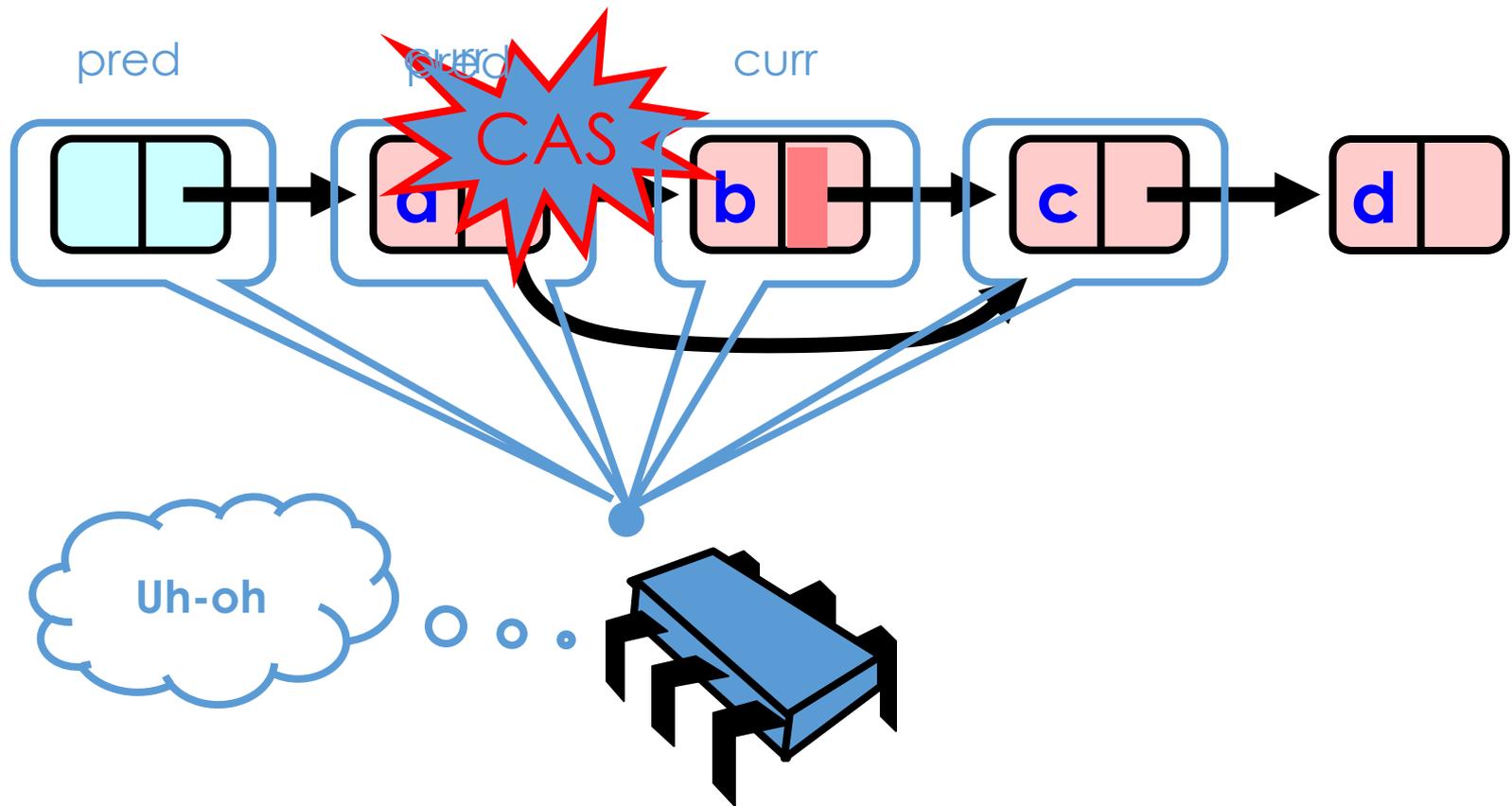
Removing a Node



Traversing the List

- Q: what do you do when you find a “logically” deleted node in your path?
- A: finish the job.
 - CAS the predecessor’s next field
 - Proceed (repeat as needed)

Lock-Free Traversal (only Add and Remove)



The END
