

Ficha 2 - Redes de Computadores

1)

$$F = 1000 \text{ KB}$$

$$RTT = 100 \text{ ms}$$

$$\text{handshake} = 1 \text{ RTT}$$

$$\text{pacote} = 1 \text{ KB}$$

b) $V_T = 1 \text{ Mbps}$

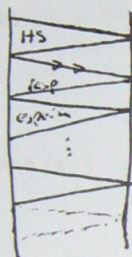
$$T_t = HS + 1000(T_{t,p} + T_{p,p}) - RTT(\text{ultimo}) + T_p$$

$$HS = 1 \text{ RTT} = 100 \text{ ms}$$

$$T_{t,p} = 8 \text{ ms} \quad T_{p,p} = RTT/2$$

$$T_t = 100 + 1000(8 + 50) + 100 + 50$$

$$= 108050 \text{ ms}$$



a) $V_T = 1 \text{ Mbps}$

$$1 \text{ pacote} \rightarrow 1 \text{ KB}$$

$$x \text{ pac.} \rightarrow 1000 \text{ KB}$$

$$x = 1000 \text{ pacotes de } 1 \text{ KB}$$

$$T_{t,p} = \frac{8000}{10^6} = 8 \text{ ms}$$

$$1 \text{ KB} = 2^{10} \text{ bytes} = 1024 \text{ bytes}$$

$$1 \text{ byte} \rightarrow 8 \text{ bits}$$

$$1024 \text{ bytes} \rightarrow 8192 \text{ bits} (= 1 \text{ KB})$$

$$T_{\text{total}} = RTT_{\text{handshake}} + T_{t,\text{pacotes}} + T_{p,p}$$

$$= 100 \text{ ms} + \frac{1000 \times 8 \text{ ms}}{16} + \frac{RTT}{2} =$$

$$= 100 \text{ ms} + 50 \text{ ms} + 50 \text{ ms} = 8,5 \text{ s}$$

d)

RTT	1°	2°	3°	N
envia	$2^0 = 1$	$2^1 = 2$	$2^2 = 4$	2^{N-1}
total env.	$2^0 - 1$	$2^2 - 1$	$2^3 - 1$	$2^N - 1$

$$T_t = 1 \text{ RTT} + N \times RTT - \frac{RTT}{2}$$

$$2^N - 1 \geq 1000 \text{ pacotes}$$

$$2^N \geq 1001 \text{ pacotes}$$

$$N \geq 9,9$$

$$T_t = 100 + 10 \times 100 - 50 = 1050 \text{ ms} = 1,05 \text{ s}$$

c) $1 \text{ RTT} \rightarrow 20 \text{ pac.}$

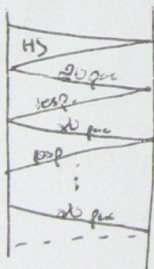
$$x \rightarrow 1000 \text{ pac.}$$

$$x = 50 \times RTT$$

$$T_t = 1 \text{ RTT} + 50 \text{ RTT} - \frac{RTT}{2}$$

$$= 100 + 5000 - 50 =$$

$$= 5050 \text{ ms} = 5,05 \text{ s}$$



1.

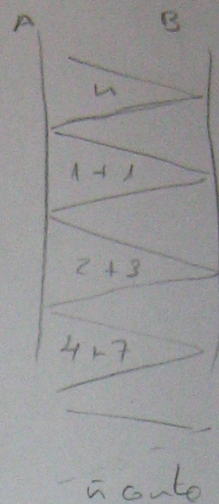
d)

$$F = 1000 \text{ KB}$$

$$P = 1 \text{ KB} \rightarrow 1000 \text{ packets}$$

$$RTT = 100 \text{ ms}$$

RTT	1^0	2^0	3^0	4^0
curia	$1 = 2^0$	$2 = 2^1$	$4 = 2^2$	2^{n-1}
total				
curved	$1 = 2^1 - 1$	$3 = 2^2 - 1$	$7 = 2^3 - 1$	$2^n - 1$



$$T = hS + N \times RTT - \frac{RTT}{2}$$

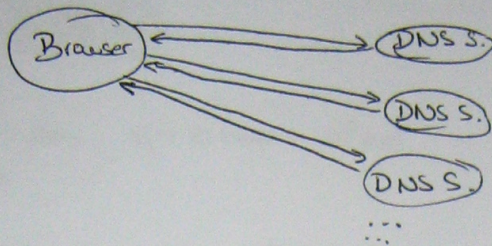
$$2^n - 1 \geq 1000 \text{ packets} \Rightarrow 2^n \geq 1001 \Rightarrow N \geq \log_2 1001 \Rightarrow N \geq 9.9$$

$$N = 10$$

$$T = 10RTT + 10RTT - \frac{RTT}{2} = 11RTT - \frac{RTT}{2}$$

$$11 \times 100 - 50 = 1100 - 50 = 1050 \text{ ms} = 1,050 \text{ s}$$

3 ④



$T_{\text{Pedido HTTP}} = \text{Handshake} + \text{pedido/resp.} = 2RTT_0$

$T_{\text{interação DNS}} = \sum_{i=1}^n RTT_i$

$T_{\text{total}} = 2RTT_0 + \sum_{i=1}^n RTT_i$

⑤

↑
DNS

1 pág + 3 obj

a) HTTP/1.0 sem conex. paralelas

$T_t = \text{DNS} + \underbrace{[1RTT + 1RTT]}_{\text{pedido página}} + 3 \times \underbrace{[1RTT + 1RTT]}_{\text{pedido 1 objeto}} = \text{DNS} + 8RTT$

Labels: abre, pede, abre, pede

b) HTTP/1.0 com conex. paralelas

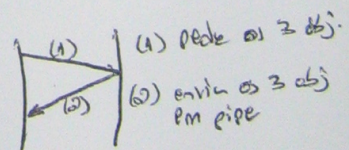
$T_t = \text{DNS} + \text{Ped. Pág} + \underbrace{3 \text{ conexões paralelas}}_{\text{cada uma } [1RTT + 1RTT]}$

$= \text{DNS} + 2RTT_0 + 2RTT_0 = \text{DNS} + 4RTT_0$

c) HTTP/1.1 com pipelining

$T_t = \text{DNS} + \text{Ped. Pág} + \text{Ped obj (3)}$

$= \text{DNS} + \underbrace{2RTT_0}_{\text{abrir + pedido}} + \underbrace{RTT_0}_{\text{conexão esta aberta}}$



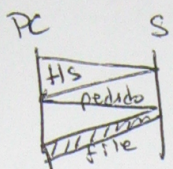
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$$RTT = 2ms \quad V_t = 10 \text{ Mbps} = 10^7 \text{ bps}$$

~~Página~~

a) $P_{\text{aj}} = 10000 \text{ bits} = 10^4 \text{ bits}$
HTTP/1.0

$$T_{\text{baj}} = \frac{10^4}{10^7} = 1ms$$



$$T_t = 2RTT + T_b = \cancel{2 \times 2ms + 1ms} = 3ms$$

$$= 2 \times 2ms + 1ms = 5ms$$

b) $RTT_{\text{ext}} = 200ms$

hitratio = 50%

$$T = 50\% \cdot [2 \times 200 + 1] +$$

$$+ 50\% \cdot [(2 \times 200 + 1) + (2 \times 200 + 1)]$$

$$= 205,5 \text{ ms}$$

