begin transaction	
update fármacos	
set NomeF= 'aspirina'	
where NumF=1;	
	update fármacos
	set NomeF= 'benuron'
	where NumF=2;
	update fármacos
	set NomeF= 'aspirina'
	where NumF=1;
update fármacos	
set NomeF= 'benuron'	
where NumF=2;	

1b) The following is a schedule for which one gets different results in the different isolation level modes. Note that for SERIALIZABLE ISOLATION LEVEL there are two possible alternative sequences <T1,T2> or <T2,T1>

T1	T2	Uncom	Commi	SR	SR
		mitted	tted	T1,T2	T2,T1
		Read	Read		
begin transaction					
select * from fármacos		empty	empty	empty	2
insert into fármacos					
values (1,'aspirina',);					
	begin transaction				
	select * from	1	empty	1	empty
	fármacos;				
	insert into fármacos				
	values (2,benuron,)				
	select * from	1	2	1	2
	fármacos	2		2	
	commit;				
select * from fármacos		1	1	1	1
		2	2		2
commit;					

1c) See slide 261 in the slides aula8concurrency.

1d) With timestamp-based protocol

TS(t1)=1	TS(t2)=2	A(R-ts,W-ts)=(0,0)	B(R-ts,W-ts)=(0,0)
read(A)		(1,0)	
	write(A)	(1,2)	
write(B)			(0,1)
	commit;		
commit;			

With 2-phase locking protocol

lock-s(A)	
read(A)	
	lock-x(A) SUSPENDS HERE !!!
	write(A)
lock-x(B)	
write(B)	
	unlock(A)
	commit;
unlock(A)	
unlock(B)	
commit;	

1e) Assuming that doctor 7 does not exists previously, the following transaction would fail if checking of foreign key constraints are not deferred to the end of the transaction, since all doctors must have at least one appointment, and all appointments are done by a doctor.

begin transaction; insert into medicos(7,'Cristiano',...); insert into consultas(1000,7,'José',...); commit;

1f) Not taught.